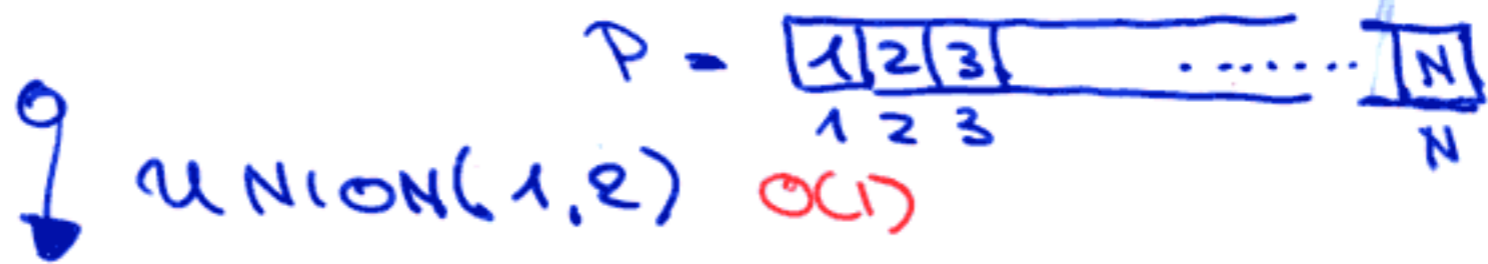
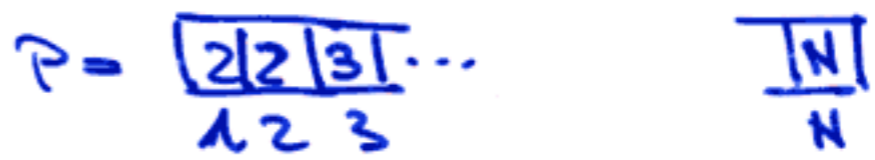


1. (1) (2) (3) (4) ... (N)



2. (2) (3) ... (N)



UNION(2,3) $\alpha(1)$

3. (3) (2) (1)

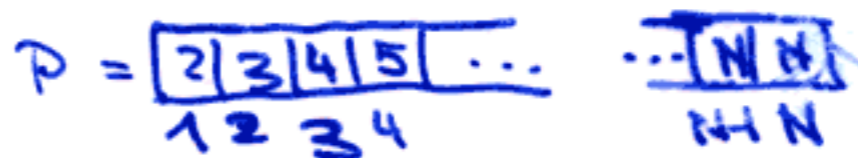


UNION(3,4)

4. (4) (2) (1)

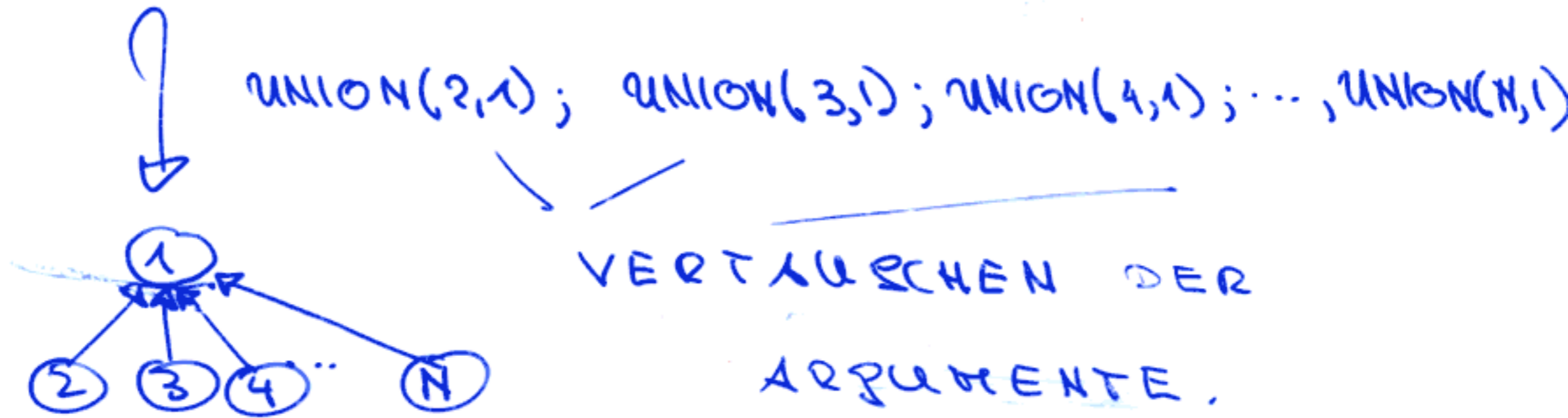


5. (N) (N-1) (2) (1)



FIND(x) IN $\Omega(N)$

① ② ③ ④ ... (N-1) (N)



FIND(i) IN $O(1)$.

⇒ UNION BY SIZE; KLEINERE ANZAHL NACH UNTEN.

UNION BY SIZE:

(1,1)

(2,1)

(3,1)

(N,1)

UNION(2,1)

(1,2)



(3,1)

(N,1)

UNION(1,3)



(4,1)

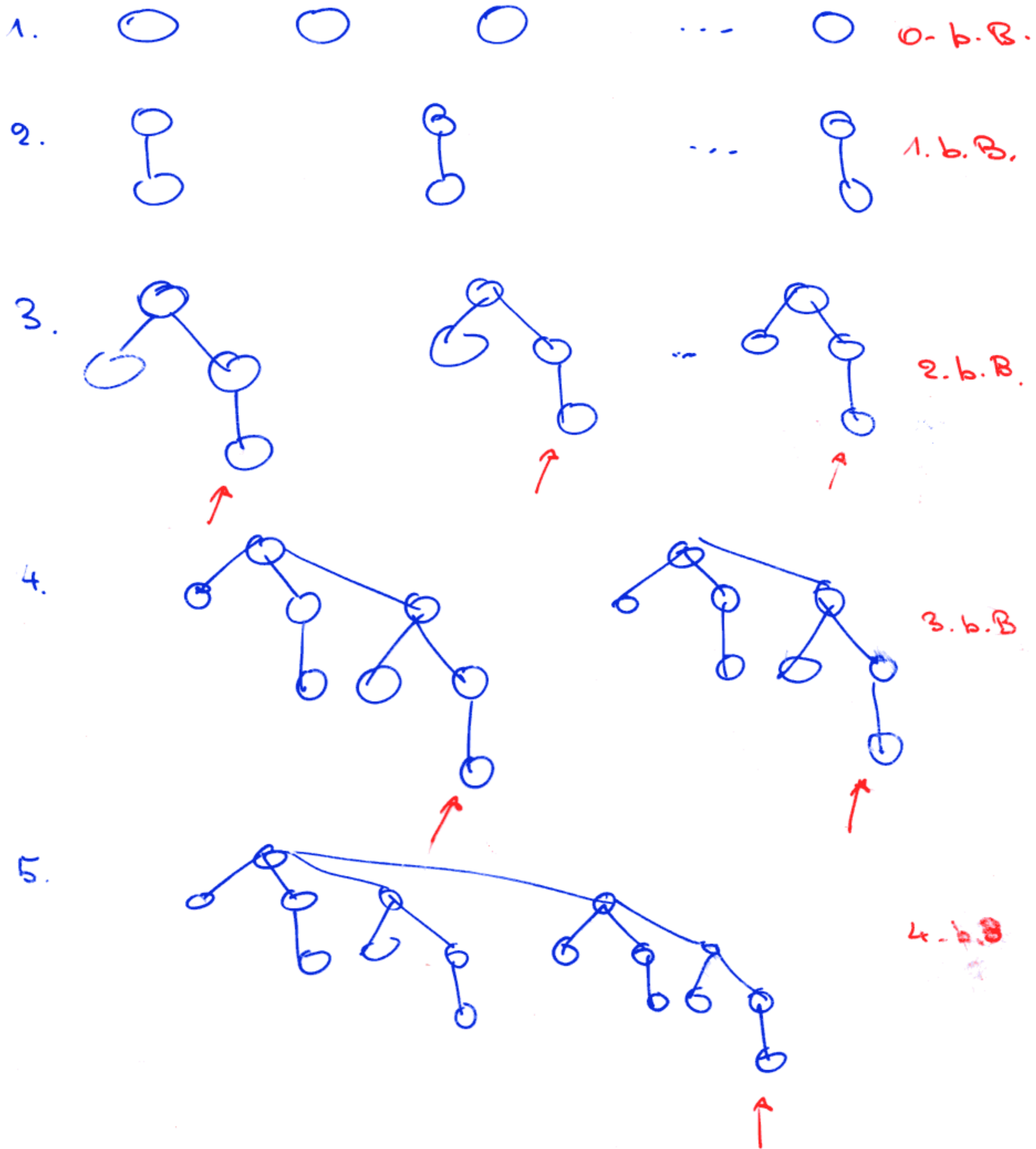
(N,1)

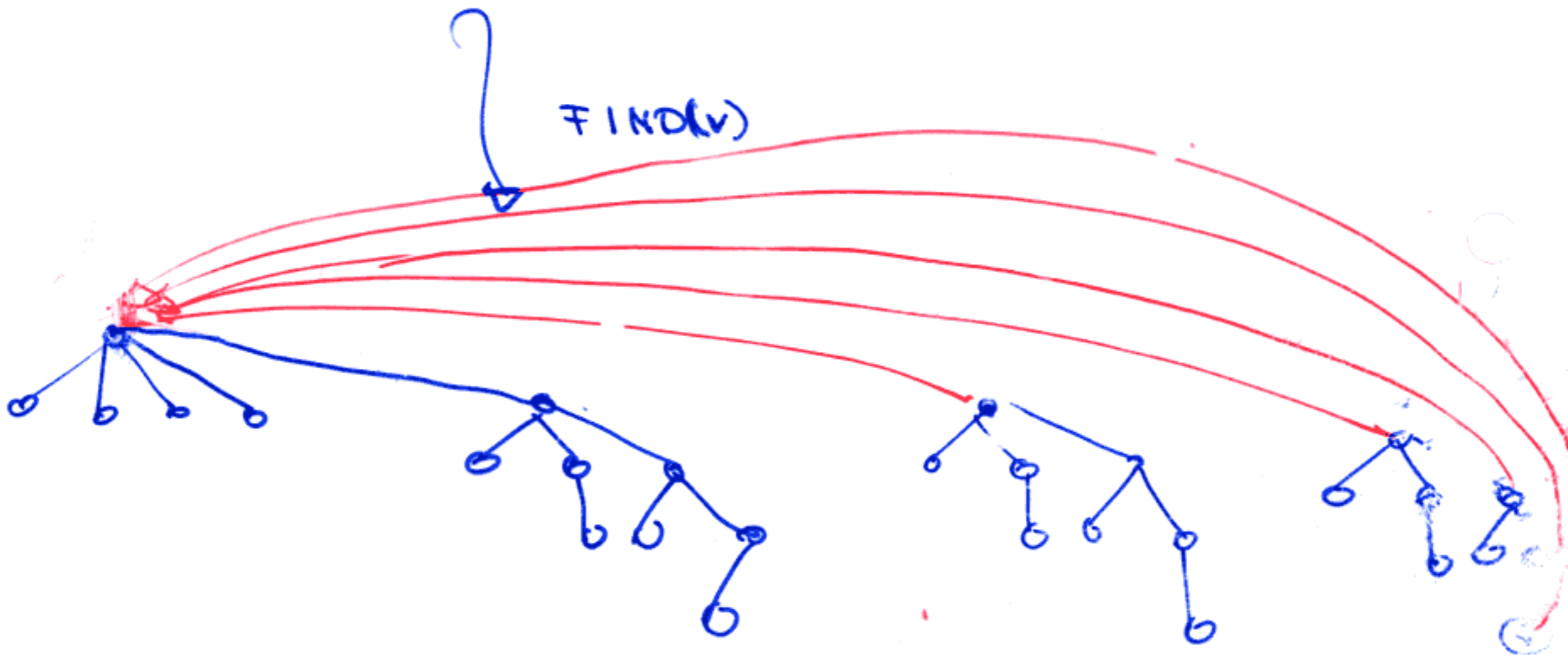
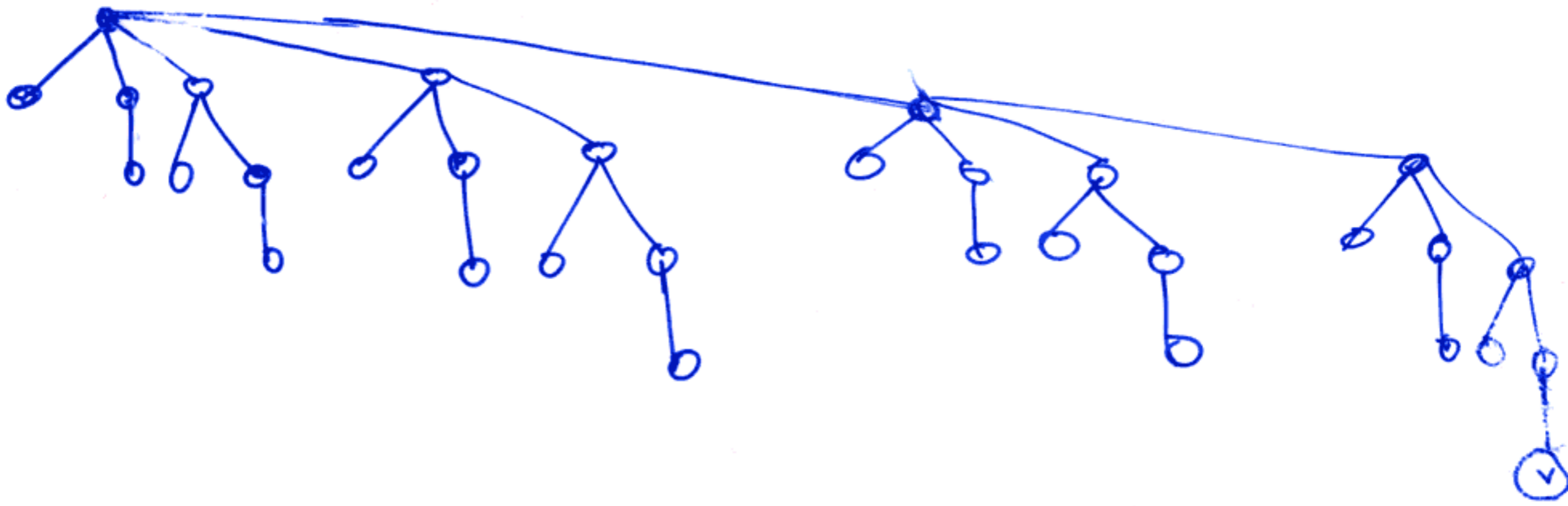
⋮

UNION(1,N)



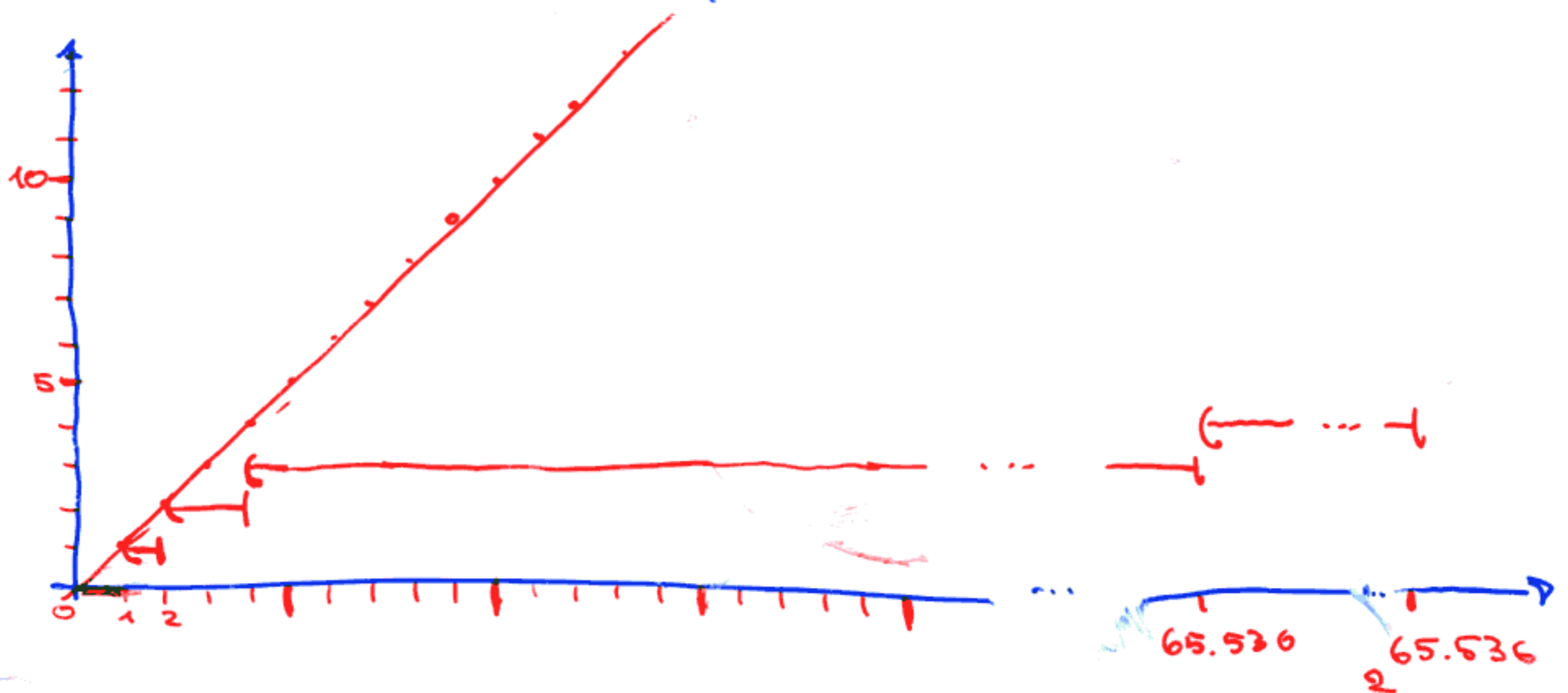
TIEFE BÄUME DURCH UNION-BY-SIZE ?





IN ALLGEMEINEN $\Omega(\log n)$ UND $O(\log n)$.

FUNKTION \log^* :



$$\log^* 1 = 0$$

$$\log^* 2 = 1$$

"FAST" KONSTANT.

$$\log^* 3 = 2$$

$$\log^* 4 = 1 + \log^* 2 = 2$$

$$\log^* 5 = 3$$

$$\log^* 8 = \log^* 3 + 1 = 3$$

$$\log^* 16 = \log^* 4 + 1 = 3$$

$$\log^* 2^{16} = 4$$

$$\log^* (2^{2^{16}}) = 5$$

$$2^{16} = 65.536$$

$$\log^* m = \text{die kleinste } s \mid \log^{(s)}(m) \leq 1$$

$$\log^{(s)}(m) = \underbrace{\log(\log(\dots(\log(m))\dots))}_{s\text{-MAL}}$$

s-MAL.

MIT UNION-BY-SIZE
UND WEGKOMPRESSION

MIT UNION-BY-SIZE
OHNE WEGKOMPRESSION

a b c d e f g h i j $\mathcal{S}_0 = \mathcal{S}'_0$...

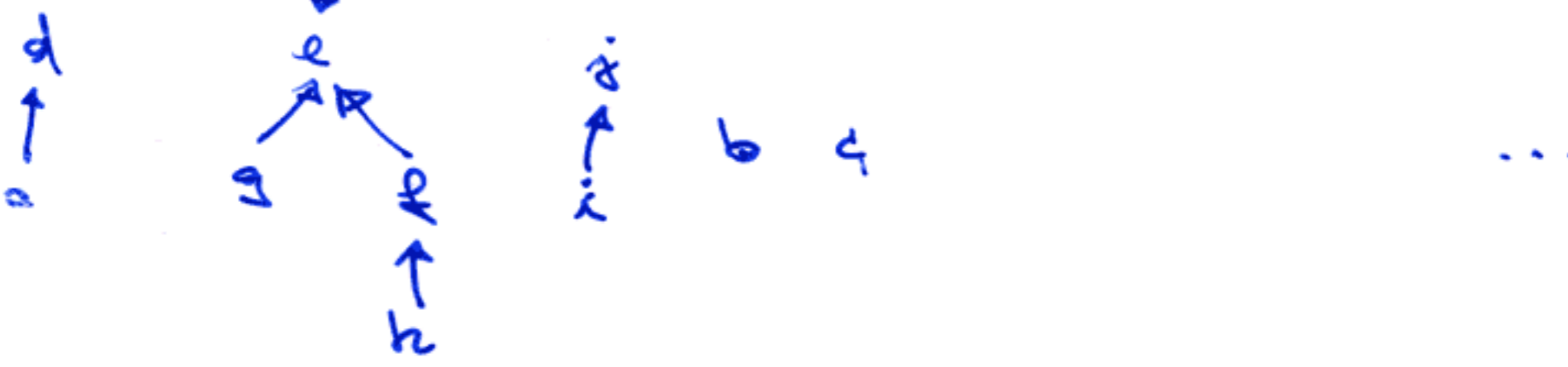
UNION(a,d); UNION(h,p); UNION(i,j)



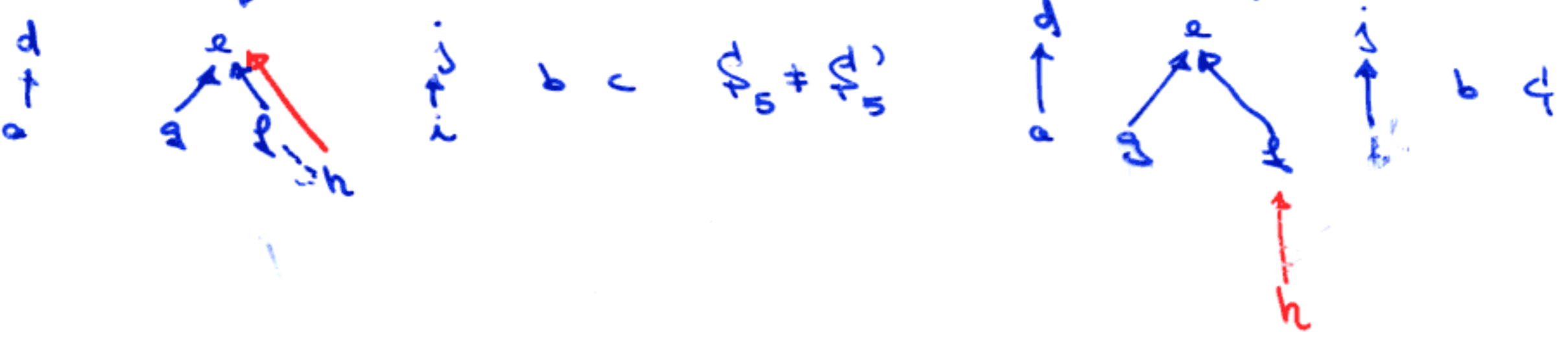
FIND(h)



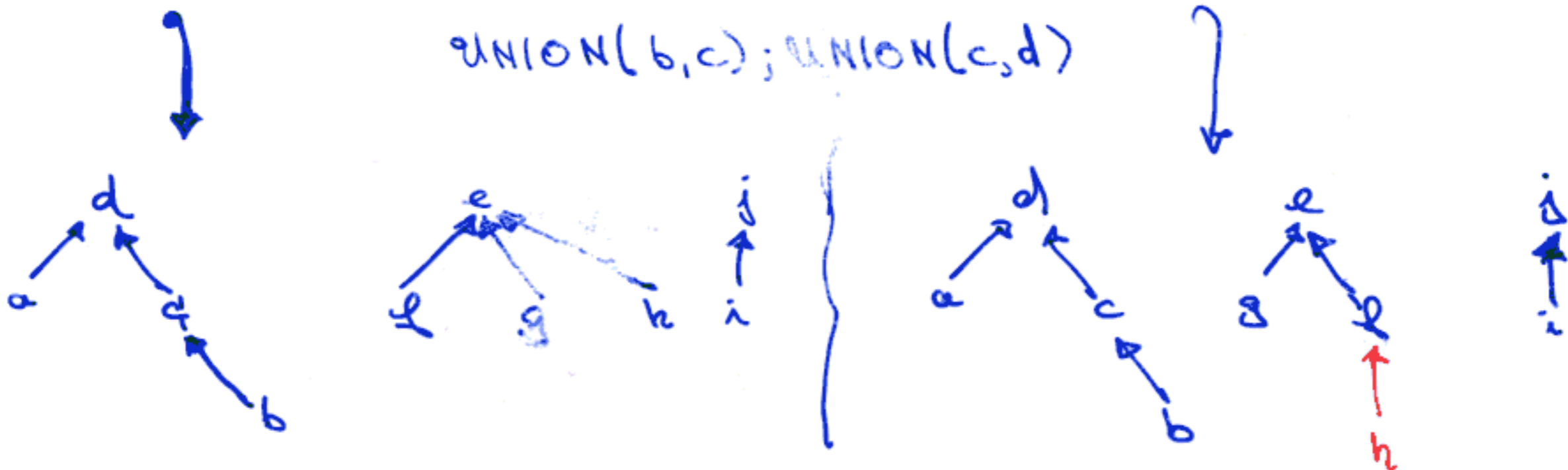
UNION(g,e); UNION(f,e)



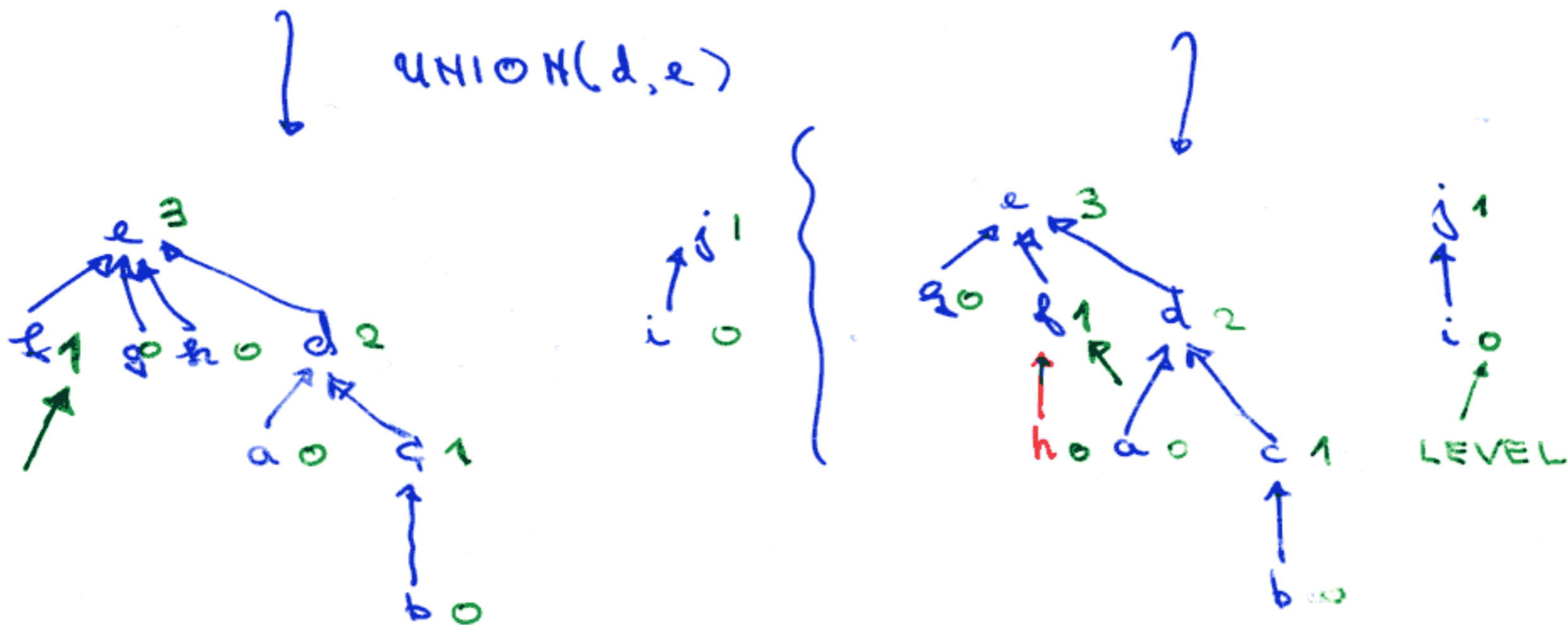
FIND(h)



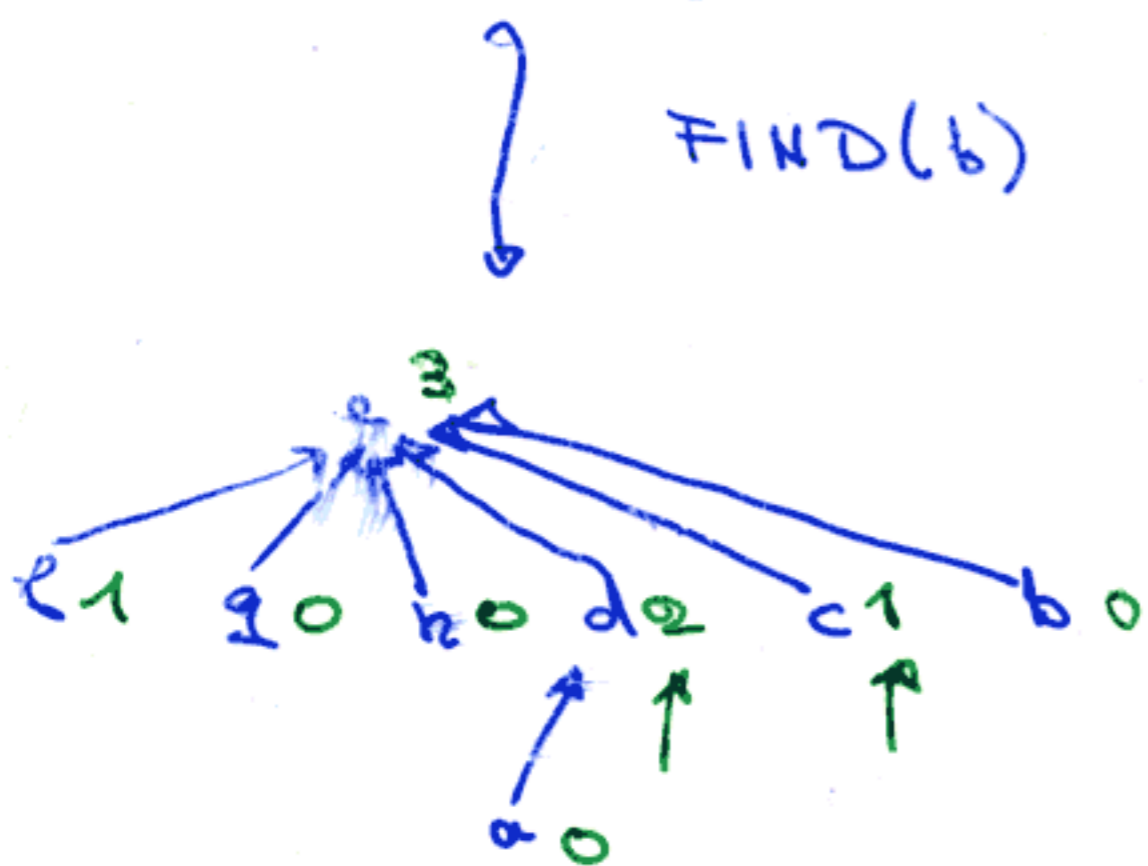
UNION(b,c); UNION(c,d)



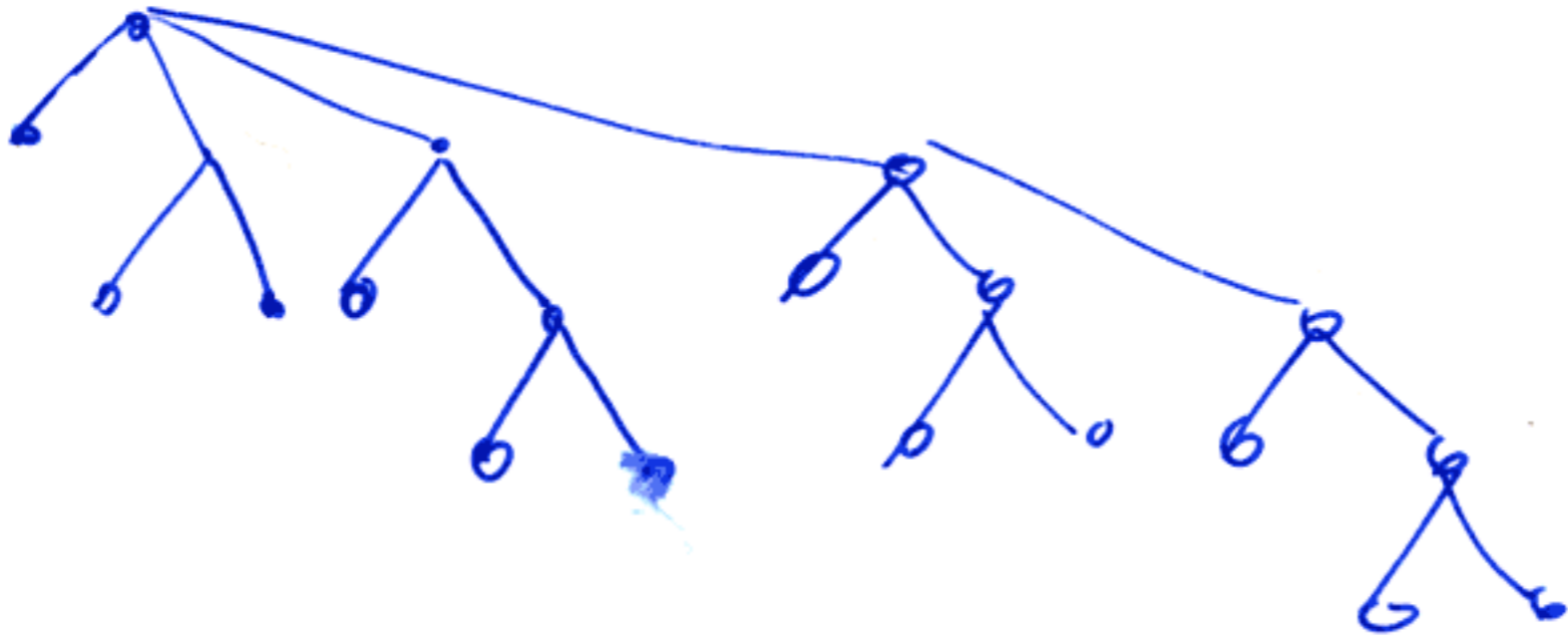
UNION(d,e)



FIND(b)



ΔN S'_m :

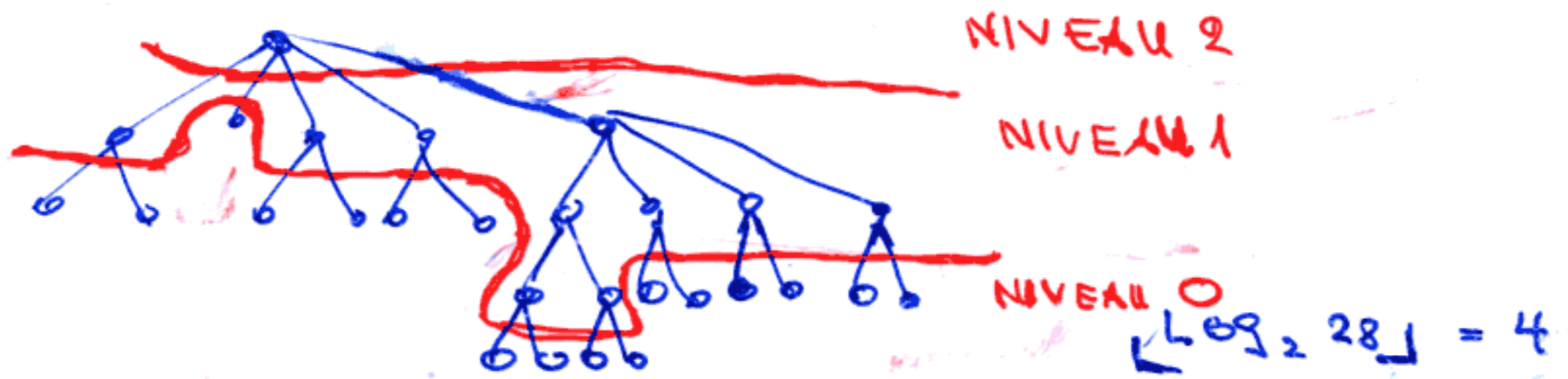


$$n = 20,$$

$$\log_2 n \downarrow = 4, \quad 4 + 2 = 6$$

$$A_0 = 0, \quad A_1 = 2, \quad A_2 = 4$$

ρ_m AUS EINEM BAUM:



$$A_0 = 0, A_1 = 1, A_2 = 4, A_3 = 5, k = 2$$



$$A_0 = 0, A_1 = 2, A_2 = 3, A_3 = 4, A_4 = 5, k = 3$$

BEACHTET: MÖGLICHE ρ_m :

